

**1. Intro.** We generate clauses that are satisfiable if and only if there's a sorting network on  $n$  lines for which (a) certain initial comparators are specified, and (b) all subsequent comparators can be done in  $t$  parallel steps.

We assume that  $n < 16$ , so that all lines can be identified by a single hexadecimal digit. Furthermore  $t$  is a single digit number, because 16 elements can be sorted in 9 rounds.

There are variables  $ijCt$ , denoting the existence of comparator  $[i:j]$  at time  $t$ ;  $kjBt$ , denoting internal nodes to check for interference/use of lines at time  $t$ ; and  $xiVt$ , denoting the value on line  $i$  at time  $t$  when the input is  $x$ . (Here  $x$  is given as four hexadecimal digits.)

```
#define maxr 20 /* at most this many initial comparators */
```

```
#include <stdio.h>
```

```
#include <stdlib.h>
```

```
int n, tt, ii, jj; /* command-line parameters */
```

```
int mask[maxr + maxr], m[16];
```

```
int leaf[32];
```

```
char needed[1 << 16];
```

```
main(int argc, char *argv[])
```

```
{
```

```
    register int i, is, j, js, k, l, r, t, x, y, z;
```

```
    <Process the command line 2>;
```

```
    <Generate the B and C clauses 3>;
```

```
    <Generate the V clauses 4>;
```

```
    <Generate the unit clauses 10>;
```

```
}
```

**2.** <Process the command line 2>  $\equiv$

```
if (argc  $\equiv$  1  $\vee$  (argc & 1)  $\equiv$  0  $\vee$  argc > maxr + maxr + 3  $\vee$ 
    sscanf(argv[1], "%d", &n)  $\neq$  1  $\vee$  sscanf(argv[2], "%d", &tt)  $\neq$  1) {
    fprintf(stderr, "Usage: %s %s %s %s\n", argv[0]);
    exit(-1);
}
```

```
if (n < 3  $\vee$  n > 15) {
    fprintf(stderr, "Sorry, %s must be between 3 and 15!\n");
    exit(-2);
}
```

```
printf("~ %s %s %s\n", n, tt);
for (j = 3; argv[j]; j++) printf("%s", argv[j]);
printf("\n");
```

This code is used in section 1.

3. There are  $n - 1$  potential comparators that use line  $k$ , namely  $[1:k]$ ,  $\dots$ ,  $[k-1:k]$ ,  $[k:k+1]$ ,  $\dots$ ,  $[k:n]$ . Place these at nodes  $n - 1$  through  $2n - 3$  of the complete binary tree with  $n - 1$  leaves. The  $n - 2$  internal nodes of this tree represent the condition that exactly one leaf in their subtree is true.

⟨Generate the B and C clauses 3⟩ ≡

```

for ( $t = 1$ ;  $t \leq tt$ ;  $t++$ )
  for ( $k = 1$ ;  $k \leq n$ ;  $k++$ ) {
    for ( $i = n - 1, j = 1$ ;  $j \leq n$ ;  $j++$ ) {
      if ( $j < k$ )  $leaf[i++] = (j \ll 4) + k$ ;
      else if ( $j > k$ )  $leaf[i++] = (k \ll 4) + j$ ;
    }
    for ( $j = n - 2$ ;  $j$ ;  $j--$ ) {
      if ( $j + j \geq n - 1$ )  $printf$  (" $\sim\%xC\%d$ ",  $leaf[j + j]$ ,  $t$ );
      else  $printf$  (" $\sim\%x\%xB\%d$ ",  $k$ ,  $j + j$ ,  $t$ );
      if ( $j + j + 1 \geq n - 1$ )  $printf$  (" $\square\sim\%xC\%d\backslash n$ ",  $leaf[j + j + 1]$ ,  $t$ );
      else  $printf$  (" $\square\sim\%x\%xB\%d\backslash n$ ",  $k$ ,  $j + j + 1$ ,  $t$ );
      if ( $j + j \geq n - 1$ )  $printf$  (" $\sim\%xC\%d\square\%x\%xB\%d\backslash n$ ",  $leaf[j + j]$ ,  $t$ ,  $k$ ,  $j$ ,  $t$ );
      else  $printf$  (" $\sim\%x\%xB\%d\square\%x\%xB\%d\backslash n$ ",  $k$ ,  $j + j$ ,  $t$ ,  $k$ ,  $j$ ,  $t$ );
      if ( $j + j + 1 \geq n - 1$ )  $printf$  (" $\sim\%xC\%d\square\%x\%xB\%d\backslash n$ ",  $leaf[j + j + 1]$ ,  $t$ ,  $k$ ,  $j$ ,  $t$ );
      else  $printf$  (" $\sim\%x\%xB\%d\square\%x\%xB\%d\backslash n$ ",  $k$ ,  $j + j + 1$ ,  $t$ ,  $k$ ,  $j$ ,  $t$ );
       $printf$  (" $\sim\%x\%xB\%d$ ",  $k$ ,  $j$ ,  $t$ );
      if ( $j + j \geq n - 1$ )  $printf$  (" $\square\%xC\%d$ ",  $leaf[j + j]$ ,  $t$ );
      else  $printf$  (" $\square\%x\%xB\%d$ ",  $k$ ,  $j + j$ ,  $t$ );
      if ( $j + j + 1 \geq n - 1$ )  $printf$  (" $\square\%xC\%d\backslash n$ ",  $leaf[j + j + 1]$ ,  $t$ );
      else  $printf$  (" $\square\%x\%xB\%d\backslash n$ ",  $k$ ,  $j + j + 1$ ,  $t$ );
    }
  }

```

This code is used in section 1.

4. ⟨Generate the V clauses 4⟩ ≡

⟨Set up the masks 5⟩;

⟨Generate all the input cases 6⟩;

**for** ( $x = 2$ ;  $x < (1 \ll n)$ ;  $x++$ )

**if** ( $needed[x]$ ) {

**for** ( $t = 1$ ;  $t \leq tt$ ;  $t++$ ) {

**for** ( $i = 1, is = 1 \ll (n - 1)$ ;  $is$ ;  $i++$ ,  $is \gg = 1$ )

**for** ( $j = i + 1, js = is \gg 1$ ;  $js$ ;  $j++$ ,  $js \gg = 1$ ) ⟨Generate the clauses for  $i:j$  on  $x$  at time  $t$  8⟩;

**for** ( $i = 1, is = 1 \ll (n - 1)$ ;  $is$ ;  $i++$ ,  $is \gg = 1$ )

⟨Generate the clauses for untouched line  $i$  on  $x$  at time  $t$  9⟩;

}

}

This code is used in section 1.

```

5. <Set up the masks 5> ≡
  for (i = 0; argv[i + i + 3]; i++) {
    if (sscanf(argv[i + i + 3], "%d", &ii) ≠ 1 ∨ sscanf(argv[i + i + 4], "%d",
      &jj) ≠ 1 ∨ ii < 1 | ii > n | jj ≤ ii | jj > n) {
      fprintf(stderr, "Invalid comparator [%s:%s]" \n", argv[i + i + 3], argv[i + i + 4]);
      exit(-3);
    }
    mask[i] = ((1 ≪ n) ≫ ii) | ((1 ≪ n) ≫ jj);
  }
  r = i;

```

See also section 7.

This code is used in section 4.

6. (Minor trick here: A binary number  $x$  is already “sorted” if and only if  $x \& (x + 1)$  is zero.)

```

<Generate all the input cases 6> ≡
  for (x = 2; x < (1 ≪ n); x++) {
    for (y = x, i = 0; i < r; i++) {
      t = mask[i] & -mask[i];
      if ((y & mask[i]) ≡ mask[i] - t) y ⊕= mask[i];
    }
    if (y & (y + 1)) needed[y] = 1;
  }

```

This code is used in section 4.

7. There are six clauses, each of which should also include  $\bar{C}_{ij}^t$ :  $(\bar{v}_i^t \vee v_i^{t-1})$ ;  $(\bar{v}_i^t \vee v_j^{t-1})$ ;  $(v_i^t \vee \bar{v}_i^{t-1} \vee \bar{v}_j^{t-1})$ ;  $(\bar{v}_j^t \vee v_i^{t-1} \vee v_j^{t-1})$ ;  $(v_j^t \vee \bar{v}_i^{t-1})$ ;  $(v_j^t \vee \bar{v}_j^{t-1})$ . We change  $v_i^t$  to  $x_i$  when  $t = 0$  and to  $y_i$  when  $t = tt$ ; and we omit clauses that are always true.

Let  $m_i = 2^{n-i} - 1$ . The first four of these clauses are always true when  $x \leq m[i]$ , because  $v_i^t = 0$  for all  $t$  in that case. Similarly, the second clause and the last three are always true when  $(x + 1) \& m[j - 1] \equiv 0$ , because  $v_j^t = 1$  for all  $t$  in that case. These simplifications are important, because the SAT solver will not immediately discover them via unit propagation.

```

#define xi (x & is)
#define xj (x & js)
#define yi (y & is)
#define yj (y & js)

```

<Set up the masks 5> +≡

```

  for (i = 0; i ≤ n; i++) m[i] = (1 ≪ (n - i)) - 1;

```

```

8. ⟨Generate the clauses for  $i:j$  on  $x$  at time  $t$  8⟩ ≡
{
  for ( $y = x$ ;  $z = y \& (y + 1)$ ;  $y -= z \gg 1$ ) ;    /* sort  $x$  */
  if ( $x \leq m[i] \vee (t \equiv 1 \wedge xi) \vee (t \equiv tt \wedge \neg yi)$ ) ;
  else {
    printf("%x%xC%d", i, j, t);
    if ( $t \neq tt$ ) printf("□~%04x%xV%d", x, i, t);
    if ( $t \neq 1$ ) printf("□%04x%xV%d", x, i, t - 1);
    printf("\n");
  }
  if ( $x \leq m[i] \vee ((x + 1) \& m[j - 1]) \equiv 0 \vee (t \equiv 1 \wedge xj) \vee (t \equiv tt \wedge \neg yi)$ ) ;
  else {
    printf("%x%xC%d", i, j, t);
    if ( $t \neq tt$ ) printf("□~%04x%xV%d", x, i, t);
    if ( $t \neq 1$ ) printf("□%04x%xV%d", x, j, t - 1);
    printf("\n");
  }
  if ( $x \leq m[i] \vee (t \equiv 1 \wedge (\neg xi \vee \neg xj)) \vee (t \equiv tt \wedge yi)$ ) ;
  else {
    printf("%x%xC%d", i, j, t);
    if ( $t \neq tt$ ) printf("□%04x%xV%d", x, i, t);
    if ( $t \neq 1$ ) {
      printf("□~%04x%xV%d", x, i, t - 1);
      if ( $((x + 1) \& m[j - 1])$ ) printf("□~%04x%xV%d", x, j, t - 1);
    }
    printf("\n");
  }
  if ( $x \leq m[i] \vee ((x + 1) \& m[j - 1]) \equiv 0 \vee (t \equiv 1 \wedge \neg xi) \vee (t \equiv tt \wedge yj)$ ) ;
  else {
    printf("%x%xC%d", i, j, t);
    if ( $t \neq tt$ ) printf("□%04x%xV%d", x, j, t);
    if ( $t \neq 1$ ) printf("□~%04x%xV%d", x, i, t - 1);
    printf("\n");
  }
  if ( $((x + 1) \& m[j - 1]) \equiv 0 \vee (t \equiv 1 \wedge \neg xj) \vee (t \equiv tt \wedge yj)$ ) ;
  else {
    printf("%x%xC%d", i, j, t);
    if ( $t \neq tt$ ) printf("□%04x%xV%d", x, j, t);
    if ( $t \neq 1$ ) printf("□~%04x%xV%d", x, j, t - 1);
    printf("\n");
  }
  if ( $((x + 1) \& m[j - 1]) \equiv 0 \vee (t \equiv 1 \wedge (xi \vee xj)) \vee (t \equiv tt \wedge \neg yj)$ ) ;
  else {
    printf("%x%xC%d", i, j, t);
    if ( $t \neq tt$ ) printf("□~%04x%xV%d", x, j, t);
    if ( $t \neq 1$ ) {
      if ( $x > m[i]$ ) printf("□%04x%xV%d", x, i, t - 1);
      printf("□~%04x%xV%d", x, j, t - 1);
    }
    printf("\n");
  }
}
}
}

```

This code is used in section 4.

9. If  $i1Bt$  is false, we have  $v_i^t = v_i^{t-1}$ .  
 ⟨Generate the clauses for untouched line  $i$  on  $x$  at time  $t$  9⟩  $\equiv$

```

{
  if ( $x \leq m[i] \vee ((x + 1) \& m[i - 1]) \equiv 0 \vee (t \equiv 1 \wedge xi) \vee (t \equiv tt \wedge \neg yi)$ ) ;
  else {
    printf ("%x1B%d",  $i, t$ );
    if ( $t \neq tt$ ) printf (" $\square\%04x\%xV\%d$ ",  $x, i, t$ );
    if ( $t \neq 1$ ) printf (" $\square\%04x\%xV\%d$ ",  $x, i, t - 1$ );
    printf ("\n");
  }
  if ( $x \leq m[i] \vee ((x + 1) \& m[i - 1]) \equiv 0 \vee (t \equiv 1 \wedge \neg xi) \vee (t \equiv tt \wedge yi)$ ) ;
  else {
    printf ("%x1B%d",  $i, t$ );
    if ( $t \neq tt$ ) printf (" $\square\%04x\%xV\%d$ ",  $x, i, t$ );
    if ( $t \neq 1$ ) printf (" $\square\%04x\%xV\%d$ ",  $x, i, t - 1$ );
    printf ("\n");
  }
}

```

This code is used in section 4.

10. Finally, we append unit clauses to suppress comparators that are redundant.

⟨Generate the unit clauses 10⟩  $\equiv$

```

for ( $i = 0; i < r; i++$ ) {
  for ( $j = i + 1; j < r; j++$ ) {
    if ( $mask[i] \& mask[j]$ ) break;
  }
  if ( $j \equiv r$ ) {
    for ( $j = 1, t = 1 \ll (n - 1); t; j++, t \gg= 1$ ) {
      if ( $mask[i] \& t$ ) break;
    }
    for ( $k = j + 1, t \gg= 1; t; k++, t \gg= 1$ ) {
      if ( $mask[i] \& t$ ) break;
    }
    printf (" $\sim\%x\%xC1\n$ ",  $j, k$ );
  }
}

```

This code is used in section 1.

**11. Index.***argc*: 1, 2.*argv*: 1, 2, 5.*exit*: 2, 5.*fprintf*: 2, 5.*i*: 1.*ii*: 1, 5.*is*: 1, 4, 7.*j*: 1.*jj*: 1, 5.*js*: 1, 4, 7.*k*: 1.*l*: 1.*leaf*: 1, 3.*m*: 1.*main*: 1.*mask*: 1, 5, 6, 10.*maxr*: 1, 2.*n*: 1.*needed*: 1, 4, 6.*printf*: 2, 3, 8, 9, 10.*r*: 1.*sscanf*: 2, 5.*stderr*: 2, 5.*t*: 1.*tt*: 1, 2, 3, 4, 7, 8, 9.*x*: 1.*xi*: 7, 8, 9.*xj*: 7, 8.*y*: 1.*yi*: 7, 8, 9.*yj*: 7, 8.*z*: 1.

- ⟨ Generate all the input cases 6 ⟩ Used in section 4.
- ⟨ Generate the **B** and **C** clauses 3 ⟩ Used in section 1.
- ⟨ Generate the **V** clauses 4 ⟩ Used in section 1.
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# SAT-MINTIME-SORT

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