

May 19, 2018 at 02:31

**1.\* Intro.** This program generates clauses that enforce the constraint  $x_1 + \dots + x_n \leq r$ , using a method due to Olivier Bailleux and Yacine Boufkhad [*Lecture Notes in Computer Science* **2833** (2003), 108–122]. It introduces at most  $(n - 2)r$  new variables  $B_{i,j}$  for  $2 \leq i < n$  and  $1 \leq j \leq r$ , and a number of clauses that I haven't yet tried to count carefully, but it is at most  $O(nr)$ . All clauses have length 3 or less.

This version inputs a graph (specified as the third parameter), and the total number of colors (as the fourth). The output clauses will insist that at least  $r$  vertices can be colored with more than one color.

```
#define nmax 10000
#include <stdio.h>
#include <stdlib.h>
#include "gb_graph.h"
#include "gb_save.h"
int n, r, kk; /* the given parameters */
int count[nmax + nmax]; /* the number of leaves below each node */
main(int argc, char *argv[])
{
    register int i, j, k, jl, jr, t, tl, tr;
    Graph *g;
    ⟨Process the command line 2*⟩;
    ⟨Output clauses for multicolored vertices 7*⟩;
    if (r ≡ 0) ⟨Handle the trivial case directly 6⟩
    else {
        ⟨Build the complete binary tree with n leaves 3⟩;
        for (i = n - 2; i; i--) ⟨Generate the clauses for node i 4*⟩;
        ⟨Generate the clauses at the root 5⟩;
    }
}
```

```

2* <Process the command line 2* > ≡
  if (argc ≠ 5 ∨ sscanf(argv[1], "%d", &n) ≠ 1 ∨ sscanf(argv[2], "%d", &r) ≠ 1 ∨ sscanf(argv[4], "%d", &kk) ≠ 1)
  {
    fprintf(stderr, "Usage: %s %d %d %d %d\n", argv[0]);
    exit(-1);
  }
  g = restore_graph(argv[3]);
  if (¬g) {
    fprintf(stderr, "I can't input the graph '%s'!\n", argv[3]);
    exit(-2);
  }
  if (g-n ≠ n) fprintf(stderr, "Warning: The graph has %d vertices, not %d!\n", g-n, n);
  r = n - r; /* x1 + ⋯ + xn ≥ r iff  $\bar{x}_1 + \dots + \bar{x}_n \leq n - r$  */
  if (n > nmax) {
    fprintf(stderr, "Recompile me: I don't allow %d\n", nmax);
    exit(-2);
  }
  if (r < 0 ∨ r ≥ n) {
    fprintf(stderr, "Eh? r should be between 0 and n-1!\n");
    exit(-2);
  }
  printf("~sat-threshold-bb %d %d\n", n, r);

```

This code is used in section 1\*.

**3.** The tree has  $2n - 1$  nodes, with 0 as the root; the leaves start at node  $n - 1$ . Nonleaf node  $k$  has left child  $2k + 1$  and right child  $2k + 2$ . Here we simply fill the *count* array.

```

<Build the complete binary tree with  $n$  leaves 3 > ≡
  for (k = n + n - 2; k ≥ n - 1; k--) count[k] = 1;
  for (; k ≥ 0; k--) count[k] = count[k + k + 1] + count[k + k + 2];
  if (count[0] ≠ n) fprintf(stderr, "I'm totally confused.\n");

```

This code is used in section 1\*.

**4\*** If there are  $t$  leaves below node  $i$ , we introduce  $k = \min(r, t)$  variables  $B_{i+1}.j$  for  $1 \leq j \leq k$ . This variable is 1 if (but not only if) at least  $j$  of those leaf variables are true. If  $t > r$ , we also assert that no  $r + 1$  of those variables are true.

**#define**  $xbar(k)$   $printf("%s.x", (g-vertices + (k) - n + 1)-name)$

```

⟨Generate the clauses for node  $i$  4*⟩ ≡
{
   $t = count[i]$ ,  $tl = count[i + i + 1]$ ,  $tr = count[i + i + 2]$ ;
  if ( $t > r + 1$ )  $t = r + 1$ ;
  if ( $tl > r$ )  $tl = r$ ;
  if ( $tr > r$ )  $tr = r$ ;
  for ( $jl = 0$ ;  $jl \leq tl$ ;  $jl++$ )
    for ( $jr = 0$ ;  $jr \leq tr$ ;  $jr++$ )
      if ( $(jl + jr \leq t) \wedge (jl + jr) > 0$ ) {
        if ( $jl$ ) {
          if ( $i + i + 1 \geq n - 1$ )  $xbar(i + i + 1)$ ;
          else  $printf("\sim B\%d.\%d", i + i + 2, jl)$ ;
        }
        if ( $jr$ ) {
           $printf("\square")$ ;
          if ( $i + i + 2 \geq n - 1$ )  $xbar(i + i + 2)$ ;
          else  $printf("\sim B\%d.\%d", i + i + 3, jr)$ ;
        }
        if ( $jl + jr \leq r$ )  $printf("\square B\%d.\%d\n", i + 1, jl + jr)$ ;
        else  $printf("\n")$ ;
      }
}
}

```

This code is used in section 1\*.

**5.** Finally, we assert that at most  $r$  of the  $x$ 's are true, by implicitly asserting that the (nonexistent) variable  $B_{1.r+1}$  is false.

⟨Generate the clauses at the root 5⟩ ≡

```

 $tl = count[1]$ ,  $tr = count[2]$ ;
if ( $tl > r$ )  $tl = r$ ;
for ( $jl = 1$ ;  $jl \leq tl$ ;  $jl++$ ) {
   $jr = r + 1 - jl$ ;
  if ( $jr \leq tr$ ) {
    if ( $1 \geq n - 1$ )  $xbar(1)$ ;
    else  $printf("\sim B2.\%d", jl)$ ;
     $printf("\square")$ ;
    if ( $2 \geq n - 1$ )  $xbar(2)$ ;
    else  $printf("\sim B3.\%d", jr)$ ;
     $printf("\n")$ ;
  }
}
}

```

This code is used in section 1\*.

6.  $\langle$  Handle the trivial case directly 6  $\rangle \equiv$

```

{
  for ( $i = 1; i \leq n; i++$ ) {
     $xbar(n - 2 + i)$ ;
     $printf("\n")$ ;
  }
}

```

This code is used in section 1\*.

7\*  $\langle$  Output clauses for multicolored vertices 7\*  $\rangle \equiv$

```

for ( $k = 0; k < g-n; k++$ )
  for ( $i = 1; i \leq kk; i++$ ) {
    for ( $j = 1; j \leq kk; j++$ )
      if ( $j \neq i$ )  $printf("\_s.\%d", (g-vertices + k)-name, j)$ ;
       $printf("\_~s.x\n", (g-vertices + k)-name)$ ;
  }
}

```

This code is used in section 1\*.

**8\* Index.**

The following sections were changed by the change file: 1, 2, 4, 7, 8.

*argc*: 1\*, 2\*  
*argv*: 1\*, 2\*  
*count*: 1\*, 3, 4\*, 5.  
*exit*: 2\*  
*fprintf*: 2\*, 3.  
*Graph*: 1\*  
*i*: 1\*  
*j*: 1\*  
*jl*: 1\*, 4\*, 5.  
*jr*: 1\*, 4\*, 5.  
*k*: 1\*  
*kk*: 1\*, 2\*, 7\*  
*main*: 1\*  
*n*: 1\*  
*name*: 4\*, 7\*  
*nmax*: 1\*, 2\*  
*printf*: 2\*, 4\*, 5, 6, 7\*  
*r*: 1\*  
*restore\_graph*: 2\*  
*scanf*: 2\*  
*stderr*: 2\*, 3.  
*t*: 1\*  
*tl*: 1\*, 4\*, 5.  
*tr*: 1\*, 4\*, 5.  
*vertices*: 4\*, 7\*  
*xbar*: 4\*, 5, 6.

- ⟨Build the complete binary tree with  $n$  leaves 3⟩ Used in section 1\*.
- ⟨Generate the clauses at the root 5⟩ Used in section 1\*.
- ⟨Generate the clauses for node  $i$  4\*⟩ Used in section 1\*.
- ⟨Handle the trivial case directly 6⟩ Used in section 1\*.
- ⟨Output clauses for multicolored vertices 7\*⟩ Used in section 1\*.
- ⟨Process the command line 2\*⟩ Used in section 1\*.

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